

An efficient algorithm for the provision of a
day-ahead modulation service by a load aggregator
ISGT 2013

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8th October 2013

- 1 Context
- 2 Day-ahead optimization
- 3 Heuristic algorithm
- 4 Conclusion

- 1 Context
 - Balancing service & consumption
 - Flexible loads management
 - Example of modulation
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Balancing service & consumption

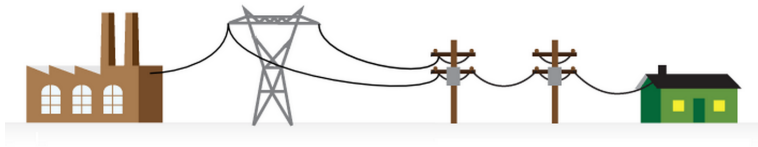
Production = Consumption: How?

Throughout the day: Balancing service

- Use flexibility of production/consumption
- ↗ needs with renewable energies

Topic of the presentation

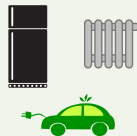
- Objective: Modulation of the consumption to offer a balancing service to the electric network.
- Coordination entity: **Aggregator**



Flexible loads management

Loads

- 1 Flexible:
Heat pump, Electric heater, Electric car, ...
- 2 Non-flexible:
TV, light, computer, micro-wave, ...



Aggregator

- Controls the consumptions of flexible loads in its portfolio.
- Provides services to the TSO.

Service proposed

- Upward/downward modulation of the consumption.
- Option available once a day.
- Activation period unknown one day ahead.

Example of modulation

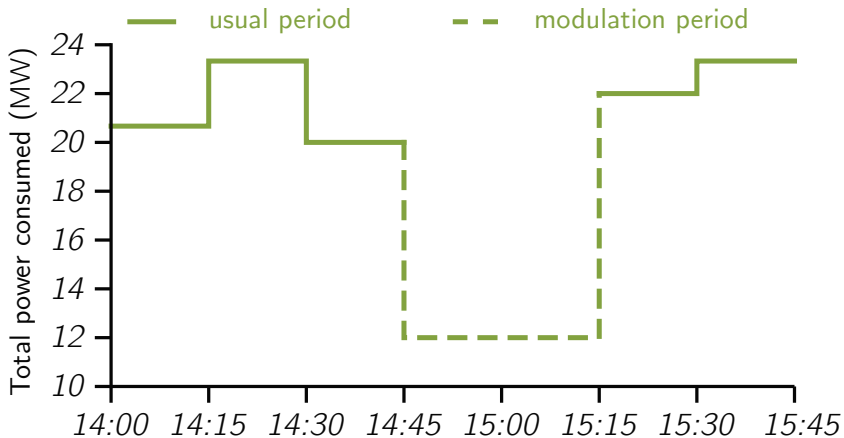


Figure: 8 MW modulation for 30 minutes.

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 - Overview
 - Generic load model
 - Mixed Integer Linear Programming
- 3 Heuristic algorithm
- 4 Conclusion

Overview

Problem considered

How to **know** and **maximize** the **quantity** of load **modulation** that can be provided given a portfolio of loads?

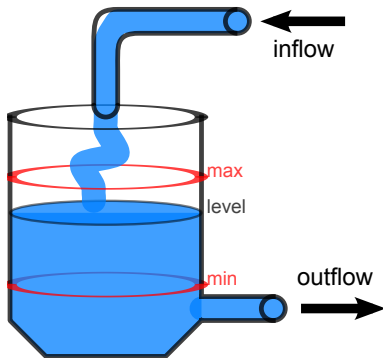
Optimization problem

max Upward and downward modulation quantities

subject to

- Feasible delivery of the modulation at any period.
ex: Can I deliver at 3 p.m?
- Load constraints
ex: The room should stay between 18°C and 22°C.

Generic load model

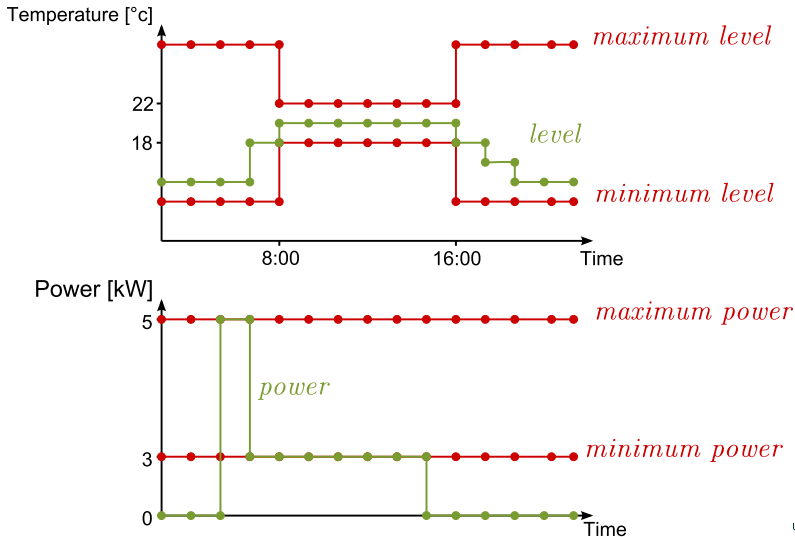


- Various type of load, one model
- Deterministic
- Related to storage
- Simple
- Level inflow = $f(\text{power})$

Level evolution

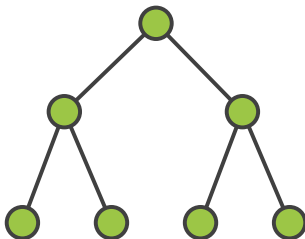
$$\Delta \text{level} = \underbrace{a_{t-\delta} \times \text{power}_{t-\delta} + b_{t-\delta} \times \text{on}_{t-\delta}}_{\text{inflow}} - \underbrace{\phi_{\text{out}_t}}_{\text{outflow}}$$

Example: Modeling of an electric heater.



Mixed Integer Linear Programming

- Mixed Integer Linear Program
 - Number of binary variables:
 - One by load and by period for each possibility of modulation call.
- ⇒ $M \times T^2$: With M loads and T periods.
- ⇒ Intractable for real instances



- 1 Context
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 - The algorithm
 - Results
- 4 Conclusion

The algorithm

Summary

- 1 Check that a feasible consumption pattern exists for each load
 - Simple set of inequations to check
- 2 Build initial solution
 - From a target total consumption at each period,
 - dispatch this consumption to every load,
 - ensure that each level is within bounds.
- 3 Local search
 - Modify the target total consumption

Complexity

$M \cdot T \cdot \log(M) + 2 \cdot M \cdot T^3$ in the worst case.

M = number of loads and T = number of periods.

1. Check feasibility for each load

Aggregator has to ensure that at least one feasible consumption pattern exists :

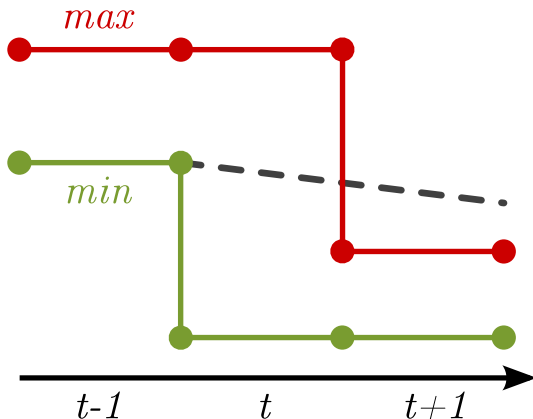


Figure: Infeasible level margins due to a too sharp decrease in the level bounds.

2. Build initial solution

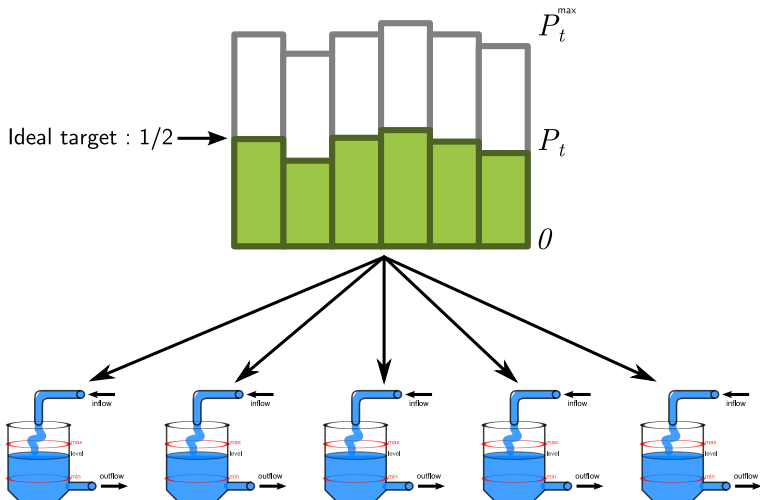


Figure: Dispatch of the reference power to every load.

3. Local search

Insight on the solution

Maximize the modulation quantity available the whole day



Maximize the **minimum** available modulation quantity.

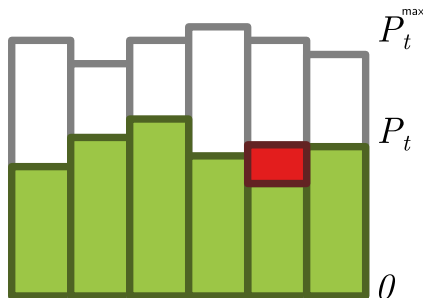


Figure: Increase of the total power reference to increase the potential of a downward modulation.

Results

- Quality :

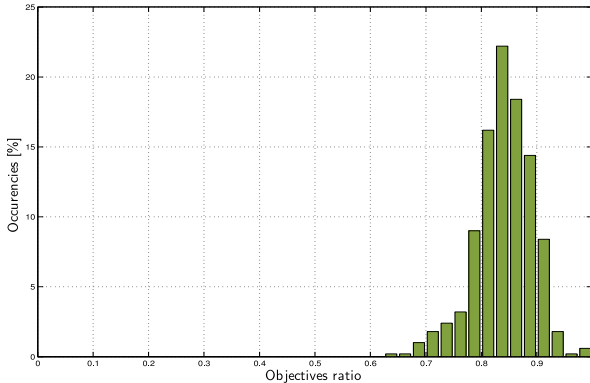


Figure: Comparison between the MIP objective value and the heuristic for 500 small instances.

- Time : 5 seconds to compare with 3 minutes for CPLEX

Conclusion

Acknowledgments :

